## Exploration of Toggle on Graphs

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### Overview

# 1 Intro

### 2 About Math Research

- How is Math Researched
- Why is Math Researched

### 3 Toggle

- Preliminary Info
- What is Toggle
- Results

• Read Existing Literature

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- Fail, try again

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- An argument is a sequence of true statements that follow one another
- Make as few assumptions as possible

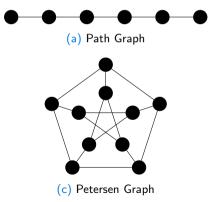
### Why do we do Math Research

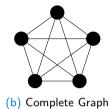
# Because its fun

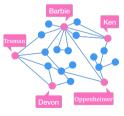
• Applications may exist in future

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- Explore the proccess of creating precise arguments









(d) Social Network

#### Definition (Impartial Two-Player Game)

An **impartial two-player game** is a game where the allowable moves only depend on the board state or position and not on which of the two players is currently moving.

### Definition (N-Position)

A two-player impartial game is in an **N-position** if that position secures a win for the **Next** player. A game is an **N-game** if the initial position is an *N*-position.

#### Definition (P-Position)

A two-player impartial game is in a **P-position** if that position secures a win for the **Previous** player. A game is a **P-game** if the initial position is a *P*-position.

#### Definition (Impartial Two-Player Game)

- Two-Player Game
- Either player may make any move

### Definition (N-Position)

• The next player to move has a winning strategy

### Definition (P-Position)

• The previous player has a winning strategy

- Impartial two player game played on a simple graph
- Players take turns turning off lights
- Last player to turn off a light wins

#### Valid Move

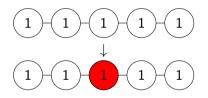
- Each light is initially on
- Toggling a light changes its flips its state and the state of its neighbors
- Only an on light may be toggled

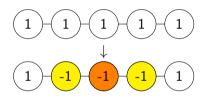
#### Valid Move

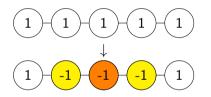
- Each light is initially on
- Toggling a light changes its flips its state and the state of its neighbors
- Only an on light may be toggled
- More lights must be off after a move

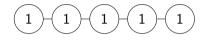
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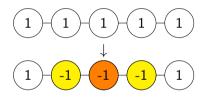
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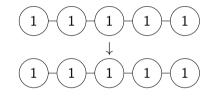


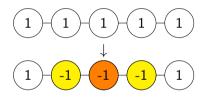


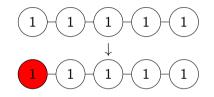


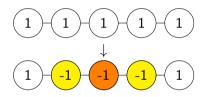


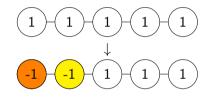


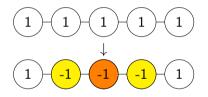


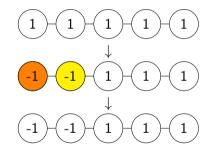


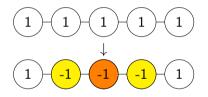


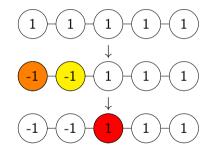


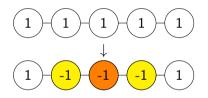


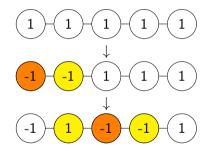












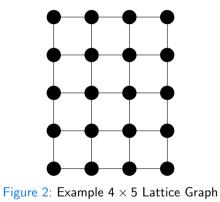
### Theorem (Sprague-Grundy Theorem)

Every impartial game under normal play is equivalent to a Nim game played on a heap of a certain size.

### Definition (Grundy Value)

The heap size of a Nim game is its **Grundy Value**. A **Grundy Value** of zero is a P-game while a nonzero **Grundy Value** is an N-game.

## Lattice Graph



### Grundy Values of $n \times m$ Latticies

$n\setminus m$	1	2	3	4	5	6
1	1	1	1	2	2	0
2	1	1	2	0	3	1
3	1	2	1	1	3	0
4	2	0	1	0	1	0
5	2	3	3	1	2	0
6	0	1	0	0	0	

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  - Memory
- Why do we care about complexity classes?
  - If two problems have the same complexity class, we can convert one into the other
  - If the algorithim for one problem is easier than another, we can convert more difficult problem into a simpler one

#### The Office



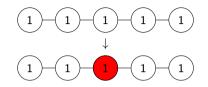
Figure 3: Scene from The Office Season 9 Episode 4

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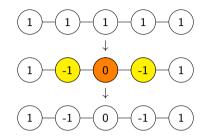
- After a light is toggled, it can not be toggled again
- The number of on lights does not need to decrease

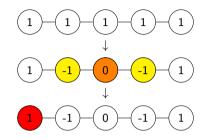
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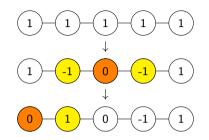
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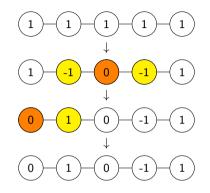


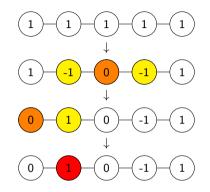
1 1 1 ↓ 0 -1 1 1 -1

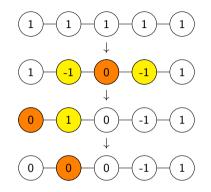


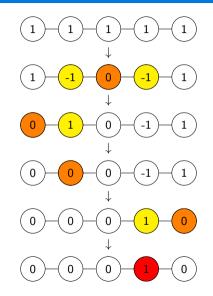








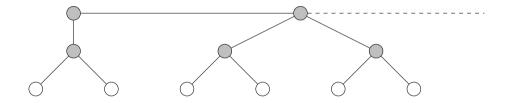




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## Charged Toggle can be modeled by Heat Toggle

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#### Heat Toggle is PSpace Complete

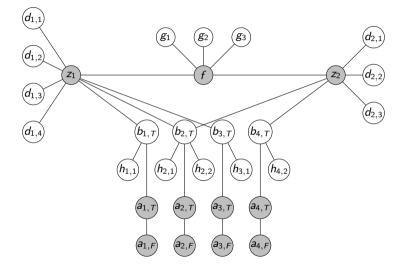


Figure 4: The expression  $(x_1 \land x_2 \land x_3) \lor (x_2 \land x_4)$  as an Heat Toggle game

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